



# Rootkits: Attacking Personal Firewalls

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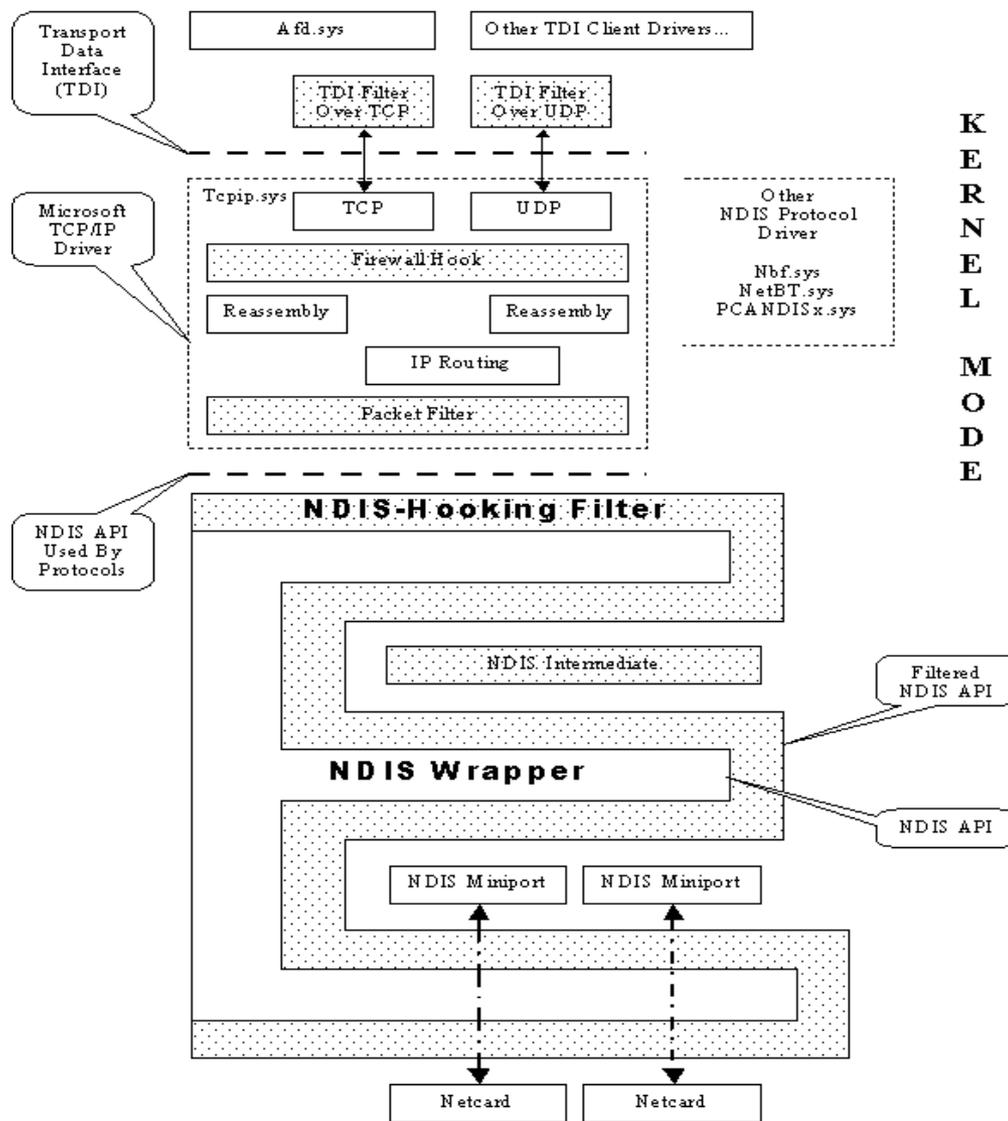
## Current personal firewalls are focused on combating usermode malware

- ▣ What about protection against rootkits?

## Overview

- ▣ i386 Windows NT+ network subsystem overview
- ▣ What malware authors usually do to cheat firewalls
- ▣ Common firewall techniques
  - ▣ Bypassing typical firewall hooks with no code patching
- ▣ Advanced firewall techniques
  - ▣ DKOM solutions to bypass modern firewalls
  - ▣ Live demo
- ▣ How to make firewalls resistant to the discussed attacks

# Windows NT Network Subsystem Overview



## Code injection into trusted process

- ▣ Malware finds trusted process and tries to inject code into it
  - ▣ Firewalls evolve to catch various types of code injections

## Prevention of firewall drivers from loading

- ▣ Rootkit registers an image load notification callback via `PsSetLoadImageNotifyRoutine()`
- ▣ The callback checks for known driver images and counteracts their loading (e.g. by patching `XOR EAX, EAX / RET 0x08` at their entry point)

- **These techniques do not actually bypass firewalls – they cheat them. They are either firewall implementation specific or take advantage of incompetence of a user (i.e. weak firewall rules exploitation)**

## TDI hooking

- ▣ Allows to implement per-process traffic monitoring and filtering connection attempts and packets of connectionless protocols made by upper-level socket interfaces
- ▣ High level TDI interfaces may be used by a firewall to simplify the detection and prevention of attacks against application layer protocols

## NDIS hooking

- ▣ Allows to implement protection against attacks targeted from data link layer (e.g. Ethernet specific attacks) to transport layer (TCP protocol attacks). TDI hooks cannot prevent data link layer attacks
- ▣ It makes possible to hook unknown protocols' traffic (for example it may be used to switch system to "network stealth" mode)

## Attaching to `\Device\Ip`, `\Device\RawIp`, `\Device\Tcp`, `\Device\Udp`

- ▣ Perform per-process traffic monitoring

## Techniques used

- ▣ Device filtering
  - ▣ Find a real device in the filter chain (lowest one)
- ▣ `DRIVER_OBJECT.MajorFunction[]` hooking
  - ▣ Perform a tunneling and find real `TCPIP.SYS` handlers

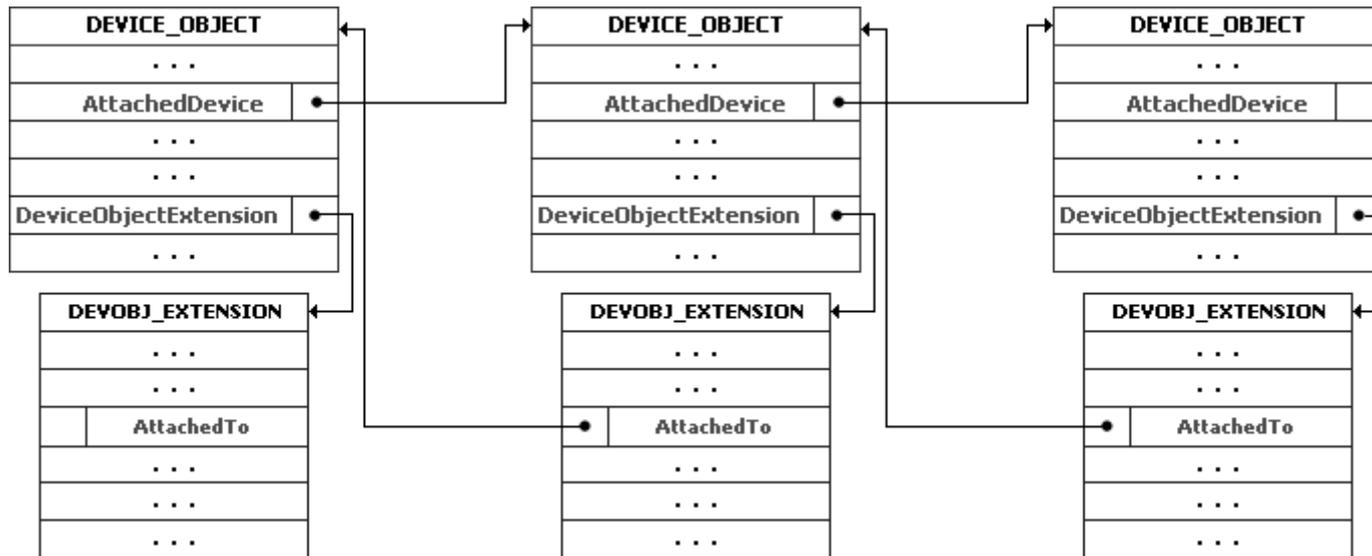
## ➤ **It's too high level to pose obstacle for a rootkit**

- ▣ It can block only rootkits that are using kernel mode sockets interface
- ▣ Greg Hogg suggested using personal TCP/IP stack in 2001

Original Device

Filter Device 1

Filter Device 2



## Walk relocation information of the `TCPIP.SYS`, store all absolute labels

- ▣ Drivers must have relocation info. Therefore, all specified MajorFunction elements must have DREFs in the driver image

## Hook `Int 01` in the IDT

- ▣ Should be done in SMP-safe way: all existing IDTs must be hooked

## Catch a thread which is going to call `IoCallDriver()` to the TDI filter device, set Trace Flag in this thread

- ▣ No code patching is required: change `pIoofCallDriver` pointer (it can be found easily – `IoofCallDriver()` is exported) like Driver Verifier or IRP Tracker does. Edgar Barbosa used `pIoofCallDriver` hooking to bypass VICE

## Trace the thread until it comes to

- ▣ one of the absolute labels of the `TCPIP.SYS`
  - ▣ Original `MajorFunction[IrpStack->MajorFunction]` found
- ▣ `IopInvalidDeviceRequest()` in the `ntoskrnl.exe`
  - ▣ This `MajorFunction` was not specified by the `TCPIP.SYS`
- ▣ the caller (`IoCallDriver()` returns)
  - ▣ TDI filter has denied something, or current IRP is pending. Wait for another `IoCallDriver()`

## Remember original `MajorFunction[ ]` value if it was found, clear TF

- ▣ Now rootkit is able to call original `MajorFunction` directly, without `IoCallDriver()`. It should adjust IRP stack locations manually

## Unhook IDT if all required `MajorFunction[ ]` entries have been found

## Hooking `NdisRegisterProtocol()`, `NdisOpenAdapter()`, `NdisDeregisterProtocol()`, `NdisCloseAdapter()`

- ▣ Catch all known (read, TCP/IP) protocol registrations
  - ▣ Patch `NDIS_PROTOCOL_CHARACTERISTICS`
- ▣ Catch all protocol bindings
  - ▣ Patch returned `NDIS_OPEN_BLOCK`

## Techniques used

- ▣ `NDIS.SYS` export table patching
- ▣ `NDIS.SYS` code patching

➤ **Still not a problem for a rootkit**

## Driver calls `NdisRegisterProtocol()`

- ▣ Firewall checks for known protocol name (usually TCPIP, RASARP and TCPIP\_WANARP) and then patches `NDIS_PROTOCOL_CHARACTERISTICS` with its own handlers
  - ▣ Some internal NDIS macros call functions by pointers from `NDIS_PROTOCOL_CHARACTERISTICS` and not from `NDIS_OPEN_BLOCK`

## Driver calls `NdisOpenAdapter()`

- ▣ Firewall calls original `NdisOpenAdapter()`, and if it succeeds, patches `(PNDIS_OPEN_BLOCK)*NdisBindingHandle` code pointers
  - ▣ `SendHandler`, `SendPacketsHandler`
  - ▣ `RequestHandler`
  - ▣ `TransferDataHandler`
  - ▣ ...

## **Rootkit may patch its own handlers over the firewall hooks in the `NDIS_OPEN_BLOCK` of a certain protocol binding**

- Used in “DeepDoor” by Joanna Rutkowska and “Peligroso” by Greg Hoglund
- This may work for simple firewalls, but more advanced ones will check their hooks for presence (subsequent `NDIS_OPEN_BLOCKS` checks) and integrity (i.e. splices/detours of their handlers)

## **How to register a protocol which will not be noticed by a NDIS-hooking firewall?**

### **Bypass firewall hooks!**

- It’s a good idea to leave hooks intact, so that firewall will notice nothing. Active antihooking may trigger the defense subsystem of the firewall

## **These hooks may be either EAT-based or direct code patches in the `NDIS.SYS`**

- EAT hooks may be defeated by finding original API addresses
- Direct code hooks may be defeated with the code pullout technique

## Load `NDIS.SYS` file image from the disk

- ▣ Assume that disk IO is not hooked. Bypassing disk IO hooks is beyond the scope of this presentation 😊

## Map image sections to appropriate virtual addresses

- ▣ This step may be skipped if we're going to translate Relative Virtual Addresses to Relative Physical Addresses using virtual section table each time we encounter a RVA. Reason: saving of memory

## Walk export table and find needed RVAs

- ▣ There's no `GetProcAddress()` equivalent in the kernel (`MmGetSystemRoutineAddress()` can be used only for `ntoskrnl` and `hal` exports)

## Apply found RVAs to the original `NDIS.SYS` image, don't rely on the import table anymore

- ▣ Make sure that API code is not hooked

## It's possible to load `NDIS.SYS` image from the disk with our own PE loader and make calls into this image

- ▣ Map `NDIS.SYS` file image sections to appropriate virtual addresses in the nonpaged memory
- ▣ All absolute pointers must be rebased to the existing `NDIS.SYS` image: we want our new hook-free code to use existing NDIS data

## Advantages

- ▣ Initialization speed: we should perform a few fairly simple operations to make things up and running
- ▣ The loaded code will always be 100% correct – it is a clone of the running NDIS
- ▣ The technique is portable: there's no need to implement different PE loader for every processor which OS supports

## Disadvantages

- ▣ Code size: we're going to use only few functions, but load the whole NDIS
- ▣ New code is identical to the original `NDIS.SYS`: a memory scanner could detect a copy

## More intelligent solution: build a sufficient NDIS code subtree

- ▣ Again, absolute pointers should be fixed to the existing NDIS image

## Advantages

- ▣ Generated code size is much smaller than the full NDIS image
- ▣ NDIS code may be mutated with any polymorphic algorithm, signatures will be broken
- ▣ If we have to perform a search for a not-exported symbol based on code XREFs or other dependencies, the searching process may be combined with code walking to improve the performance

## Disadvantages

- ▣ Initialization speed: there is a number of time-consuming operations
- ▣ It is theoretically possible that we encounter instructions that our disassembler will not be able to decode: the disassembler engine must understand as many instructions subsets as possible
- ▣ It's architecture-dependent: one has to implement code coverage and rebuilding tools for every supported processor

## PE loader maps new virtual image of the NDIS.SYS from the disk

- ▣ Don't care about relocs – they will be fixed later
- ▣ Do not use `MmCreateSection()` with a `SEC_IMAGE` allocation attribute: original section mapper (`MiCreateImageFileMap()`) may be hooked

## Entry points for the subtree are defined as RVAs of the needed APIs

- ▣ All subtrees will intersect with each other over the shared code – generated code should not be redundant

## Engine builds a code coverage map: each queued branch is being statically walked with a disassembler

- ▣ We stop on `RET`, `IRETD`, unpredictable control transfer (like `JMP reg32`) or when we come to the code that has been already analyzed. Calls and conditional jumps “fork” execution flow – they add branches to later analysis. Subtree coverage map is complete when there is no more branches left in the analysis queue

## **All contiguous regions of the covered code are copied to one chunk of memory one after another without gaps**

- ▣ Here we recalculate entry points for the addresses which were specified as the top of the original code subtrees (NDIS APIs in our case)
- ▣ This is where polymorphic methods may be applied to get rid of any static code signatures

## **Engine relinks all relative jumps and calls in the generated code**

- ▣ All relative instructions that connect non-adjacent code blocks were damaged while merging a coverage

## **Relocations are fixed to the original NDIS image**

## Registration of a dummy protocol for walking protocols list

- ▣ Will spot new protocols without hooking, thus leaving antihooking methods useless

## Periodical checks of the `NDIS_OPEN_BLOCKS` code pointers integrity

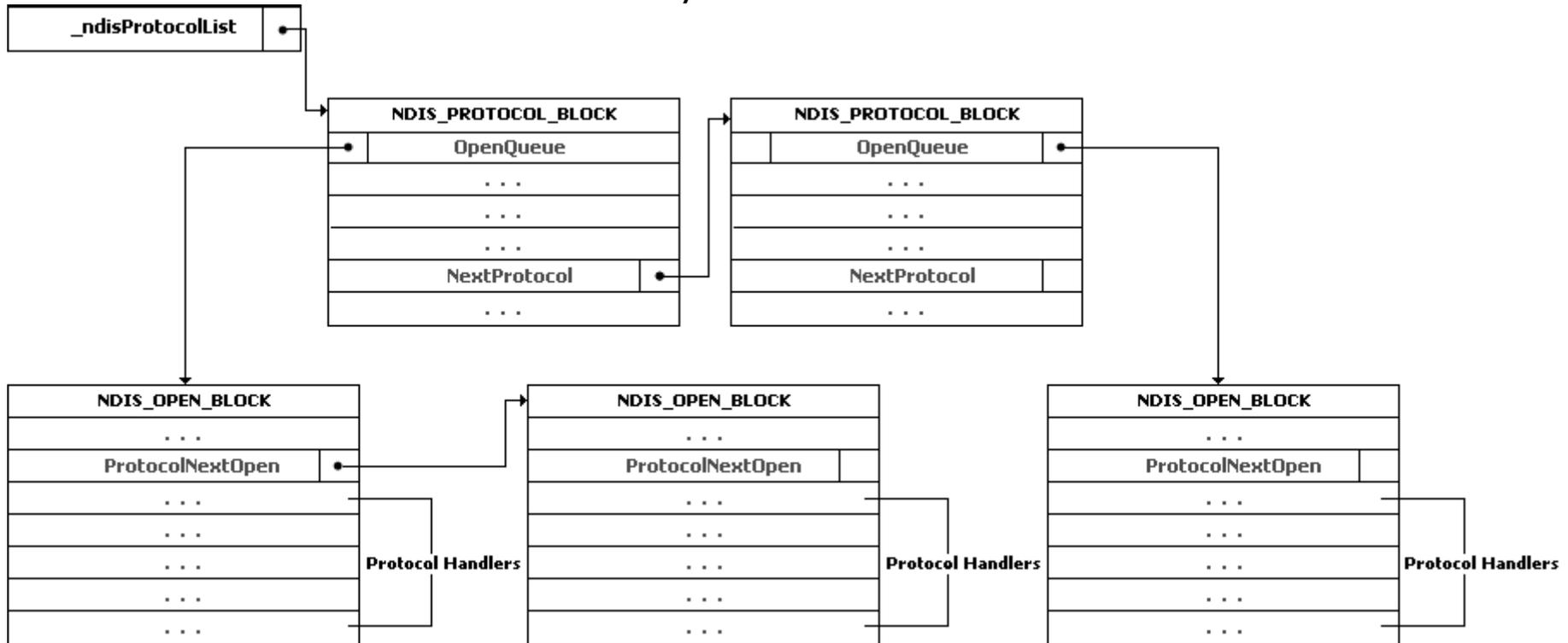
- ▣ "DeepDoor" and "Peligroso" rootkits will lose their hooks

## Anti-splice and anti-detours tricks

- ▣ Various control data is addressed in trampolines via PIC code with the help of EIP-based deltas: direct detours will change the logic of the firewall trampoline, which may lead to BSOD or rootkit compromise

**NdisRegisterProtocol () returns valid NDIS\_PROTOCOL\_BLOCK pointer which is first in the protocols list**

- Walking list this way is **dangerous!** ndisProtocolListLock must be acquired, otherwise a race condition may occur



## The right solution will be to use real `ndisProtocolList` and `ndisProtocolListLock`

- The problem: they are not exported
- `ndisProtocolList` is singly linked, so we can't walk it backwards to find a head
- To be sure that firewall is not cheating us, we again will use static analysis of the `NDIS.SYS` file
- Here's a fact: both these global variables are used by the `NdisRegisterProtocol()`

## First, enumerate all absolute pointers in the `NdisRegisterProtocol()` execution tree

- Eliminate all IAT pointers from this list

## Now check, which global variable from the list is ever used as a `PKSPIN_LOCK` by examining calls to `KfAcquireSpinLock()` and such

- From NT4 till 2003 Server there will be just one spin lock – the `ndisProtocolListLock`

## Acquire found spin lock and check other global variables – do they look like a head of a `NDIS_PROTOCOL_BLOCK` singly linked list

- Some memory forensics required!

## `NdisRegisterProtocol ()`

- ▣ Places new `NDIS_PROTOCOL_BLOCK` at the head of the `ndisProtocolList`

## `NdisDeregisterProtocol ()`

- ▣ Removes protocol from the list

## `ndisReferenceProtocolByName ()`

- ▣ `ndisCheckAdapterBindings ()`
- ▣ `ndisHandleProtocolReconfigNotification ()`
- ▣ `ndisHandleProtocolUnloadNotification ()`
- ▣ `ndisHandleProtocolBindNotification ()`
- ▣ `ndisHandleProtocolUnbindNotification ()`

## `ndisDereferenceProtocol ()`

- ▣ Decrements reference counter and frees `NDIS_PROTOCOL_BLOCK` if it reaches zero
- ▣ Does not walk `ndisProtocolList` if the protocol remains referenced

## `ndisPnPDispatch ()`

- ▣ Checks for empty `ndisProtocolList` before calling `ndisQueueBindWorkitem ()`

## `ndisCheckAdapterBindings ()`

➤ **`ndisProtocolList` is not used by the packet indication code**

**NdisOpenAdapter () updates a corresponding miniport filter database (ETH\_FILTER for ethernet in NT4/2000, X\_FILTER structure in XP+)**

- Database is selected using Miniport->MediaType value
  - Miniport->EthDB for ethernet
  - Miniport->TrDB for token ring
  - Miniport->FddiDB for fiber optic
- For ARCnet miniports ARC\_FILTER is used instead of X\_FILTER; the filter database is at Miniport->ArcDB

```
struct _X_FILTER {          // XP SP2
    /*<+0x0>*/ /*|0x4|*/ struct _X_BINDING_INFO* OpenList;
    /*<+0x4>*/ /*|0x210|*/ struct _NDIS_RW_LOCK BindListLock;
    /*<+0x214>*/ /*|0x4|*/ struct _NDIS_MINIPORT_BLOCK* Miniport;
    /*<+0x218>*/ /*|0x4|*/ unsigned int CombinedPacketFilter;
    /*<+0x21c>*/ /*|0x4|*/ unsigned int OldCombinedPacketFilter;
    /*<+0x220>*/ /*|0x4|*/ unsigned int NumOpens;
    /*<+0x224>*/ /*|0x4|*/ struct _X_BINDING_INFO* MCastSet;
    /*<+0x228>*/ /*|0x4|*/ struct _X_BINDING_INFO* SingleActiveOpen;
    /*<+0x22c>*/ /*|0x6|*/ unsigned char AdapterAddress[6];
    ...
}
```

**XNoteFilterOpenAdapter() / EthNoteFilterAdapter() attaches new ETH BINDING\_INFO / X\_BINDING\_INFO to the selected filter database**

- ▣ Current NDIS\_OPEN\_BLOCK pointer is stored there

**This way NDIS saves information about NDIS\_OPEN\_BLOCK bindings to the particular NDIS\_MINIPORT\_BLOCK**

- ▣ NDIS **does not** use ndisProtocolList to find an open binding on any network event, but firewalls do (indirectly): they get information about bindings by walking the protocol list

```
struct _X_BINDING_INFO {           // XP SP2
    /*<+0x0>*/ /*|0x4|*/ struct _X_BINDING_INFO* NextOpen;
    /*<+0x4>*/ /*|0x4|*/ struct _NDIS_OPEN_BLOCK* NdisBindingHandle;
    /*<+0x8>*/ /*|0x4|*/ void* Reserved;
    /*<+0xc>*/ /*|0x4|*/ unsigned int PacketFilters;
    /*<+0x10>*/ /*|0x4|*/ unsigned int OldPacketFilters;
    /*<+0x14>*/ /*|0x4|*/ unsigned long References;
    ...
}
```

## Ethernet packet managers

- ▣ `ethFilterDprIndicateReceivePacket()`
  - ▣ Gets `X_FILTER` pointer by looking into `PNDIS_MINIPOINT_BLOCK (Miniport->EthDB)`, which is its first parameter
- ▣ `EthFilterDprIndicateReceive()`
  - ▣ For legacy miniports only
  - ▣ Gets `ETH_FILTER/X_FILTER` pointer as the first parameter

## Managers walk `ETH_BINDING_INFO` / `X_BINDING_INFO` lists and indicate packets to the appropriate protocols

- ▣ NDIS doesn't care about `NDIS_PROTOCOL_BLOCKS` here; only `NDIS_OPEN_BLOCKS` matter (`X_BINDING_INFO.NdisBindingHandle`)

➤ Therefore, NDIS may indicate the packets to the protocol **which is not present in the `ndisProtocolList`**

## New protocol registration: Code Pullout + DKOM methods

- ▣ We should exclude our protocol from the `ndisProtocolList`: it will remain functional, but a firewall won't be able to find it using list walking
- ▣ **Approach I**: call hook-free versions of `NdisRegisterProtocol()`, `NdisOpenAdapter()` and then unlink `NDIS_PROTOCOL_BLOCK` from the list. Very similar to process hiding via `PsActiveProcessHead` elements unlinking
- ▣ **Approach II**: modify copied `NdisRegisterProtocol()` and `NdisOpenAdapter()` code trees

## Without new protocol: nothing to hide

- ▣ We should establish hooks over the existing protocols; hooking `NDIS_OPEN_BLOCKS` is too high level
- ▣ **Approach III**: hook existing `ETH_BINDING_INFO / X_BINDING_INFOS`
- ▣ **Approach IV**: register new `ETH_BINDING_INFO / X_BINDING_INFO` manually

## **Unhook `NdisRegisterProtocol ()` and `NdisOpenAdapter ()`**

### **Call these hook-free APIs to register and bind a rootkit protocol**

- ▣ It will be linked in the `ndisProtocolList`, but a firewall will not detect its registration and binding

### **Unlink returned `NDIS_PROTOCOL_BLOCK` from the list**

- ▣ We have already found `ndisProtocolList`

### **Major shortcoming**

- ▣ Firewall may detect and hook newly registered protocol before we unlink it
- **Easy to implement, but very impractical: rootkit may be compromised if a firewall has a timer which is frequent enough**

## **Modify our copies of `NdisRegisterProtocol()` and `NdisOpenAdapter()` code trees**

- ▣ Substitute all references to the original `ndisProtocolList` with references to the fake one in the generated code: both APIs will remain coherent
  - ▣ This may be done on the final step of the code generating – relocations linking
  - ▣ Fake `ndisProtocolList` may be NULL
- **Uses disassembler engine (i.e. not easily portable), requires to hook `Receive*` handlers of all other protocols bound to same adapter in order to block packets designated to our TCP/IP stack – only rootkit protocol should receive them**

## It has been shown that NDIS packet receive managers use X\_FILTER.OpenList as a head of all open bindings

- ▣ Choose random protocol binding to the specific adapter by walking its X\_FILTER.OpenList
  - ▣ Make a copy of its NDIS\_OPEN\_BLOCK (accessed via X\_BINDING\_INFO.NdisBindingHandle)
  - ▣ Patch Receive\* handlers in the copied open block
  - ▣ Substitute pointer to the original NDIS\_OPEN\_BLOCK for pointer to the patched copy
- **Very stealthy: this approach introduces only one pointer modification to a system**

## Register new `ETH_BINDING_INFO` / `X_BINDING_INFO` manually

- ▣ Create correct `NDIS_OPEN_BLOCK` without `NdisOpenAdapter()`
- ▣ Properly add new `X_BINDING_INFO` which points to faked `NDIS_OPEN_BLOCK` to the `X_FILTER.OpenList`

## Major shortcoming

- ▣ Very NDIS version dependent
- **Very hard to implement properly; different code for every supported NDIS version**

## What about sending packets?

- ▣ It's almost trivial: only `NDIS_OPEN_BLOCK` and `NDIS_MINIPORT_BLOCK` are required, and they are not hooked by a firewall

## Hook-free `NdisOpenAdapter()` may set `NDIS_OPEN_BLOCK.SendHandler` to

- ▣ `ndisMSendX()`
- ▣ `ndisMSend()`
- ▣ `ndisMWanSend()`
- ▣ `ndisMSendSG()`

## These APIs may be hooked in the `NDIS.SYS` image

- ▣ Use code pullout again – this time without any relocations updates

## We didn't register our own protocol – we hooked X\_BINDING\_INFO of an existing one

- So, its `NDIS_OPEN_BLOCK.SendHandler` and `SendPacketsHandler` may be hooked by a firewall

## In order to send packets stealthy, we should find original `ndisMSend*` functions

- By tunneling the firewall with an innocent packet: sooner or later it should be sent via call to one of NDIS packet send functions
- By searching NDIS image for not-exported symbols using XREFs or code signatures analysis
- By temporarily registering and binding a dummy protocol with aid of previously discussed methods to get original NDIS send functions pointers
  - Protocol registration and binding must not be caught by a firewall

**FireWalk rootkit:  
kernel mode FTP server over the rootkit's TCP/IP stack VS  
popular personal firewalls**

**A firewall should operate at a more privileged level than a rootkit, otherwise **it can always be bypassed****

**Since in i386 NT they both run in kernel mode, the only solution for firewall vendors is to complicate rootkits' (and their authors') life as much as possible**

- ▣ Maybe full rewrite of NDIS (with a lot of obfuscations) is a good idea – at least, there will be no symbols 😊

## Find unlinked protocols

- ▣ Walk filter databases for each miniport, get a list of `NDIS_OPEN_BLOCKS` bound to them
- ▣ Hook all found `NDIS_OPEN_BLOCKS`
- ▣ Save `NDIS_PROTOCOL_BLOCKS` associated with each `NDIS_OPEN_BLOCK`
- ▣ Walk `ndisProtocolList` and alert user about unlinked protocols

## KLISTER by Joanna Rutkowska did similar things to find processes unlinked from `PsActiveProcessHead` list

- ▣ It was bypassed too 😊

**Firewall has to take into account that rootkit may not use its `NDIS_OPEN_BLOCK.SendHandler()` or `SendPacketsHandler()` to send packets to the network**

**Rootkit may call `ndisMSend*` directly**

- ▣ However, it should find these functions first
  
- **Firewall should at least hook code of `ndisMSend*` and `ndisMWanSend*`. The nature of packet send interface does not require any special system object registration (you should register and bind a protocol in order to receive packets), and until this behavior doesn't change, firewalls will be having hard times catching sent packets**

**Joanna Rutkowska, *KLISTER***

<http://invisiblethings.org/tools/klister-0.4.zip>

**Joanna Rutkowska, *Rootkits vs. Stealth by Design Malware***

[http://invisiblethings.org/papers/rutkowska\\_bhfederal2006.ppt](http://invisiblethings.org/papers/rutkowska_bhfederal2006.ppt)

**Greg Hoggund, *NT Rootkit***

[http://www.rootkit.com/vault/hoggund/rk\\_044.zip](http://www.rootkit.com/vault/hoggund/rk_044.zip)

**Opc0de, *Bypassing VICE 2***

<http://rootkit.com/newsread.php?newsid=197>

**PCAUSA, *Windows Network Data and Packet Filtering***

<http://www.ndis.com/papers/winpktfilter.htm>

**90210, *Bypassing Klister 0.4 With No Hooks or Running a Controlled Thread Scheduler***

<http://www.rootkit.com/vault/90210/phide2.zip>

**Thank you for your time!**