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## [ ADVANCED MAC OS X PHYSICAL MEMORY ANALYSIS ]

In 2008 and 2009, companies and governments (e.g. Law Enforcement agencies) interests for Microsoft Windows physical memory grew significantly. Now it is time to talk about Mac OS X. This paper will introduce basis of Mac OS X Kernel Internals regarding management of processes, threads, files, system calls, kernel extensions and more. Moreover, we are going to details how to initialize and perform a virtual to physical translation under an x86 Mac OS X environment.

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# ADVANCED MAC OS X PHYSICAL MEMORY ANALYSIS

## INTRODUCTION

In 2008 and 2009, companies and governments (e.g. Law Enforcement agencies) interests for Microsoft Windows physical memory grew significantly. Now it is time to talk about Mac OS X. This paper introduces Mac OS X Kernel Internals regarding management of processes, threads, files, system calls, kernel extensions and more. We provide details on how to initialize and perform a virtual to physical translation under a x86 Mac OS X environment.

Physical Memory is widely known in the UNIX world as `/dev/mem`.

## MEMORY ADDRESS TRANSLATION

### QUICK TRANSLATION FORMULA

Most Operating Systems have a way to compute the kernel physical address even if you do not have the `cr3` register value which is used as Directory Table Base for virtual to physical address translation. If you want to have more detailed information on this, please refer to [\*Intel64 and IA-32 Architectures Software Developer's Manual: Volume 3A System Programming Guide\*](#).<sup>1</sup>

By kernel physical addresses, I mean the kernel image (`__DATA` & `__CODE` sections) physical address. Both contain important information and variables we need. For instance, to reconstruct the kernel address space we need to be able to use Smart Translation Formula which requires variables we can retrieve using Quick Translation Formula. As I said above, with Quick Translation Formula we can only access to `__DATA` and `__CODE` sections of the kernel image and not to allocated buffers.

Here is a summary of some operating systems with their corresponding formula to translate from Kernel Virtual Address (KVA) to Kernel Physical Address (KPA).

Operating System	Quick translation formula
x86 Linux	$KPA = KVA - 0xC0000000$
PlayStation 3 Linux	$KPA = KVA - 0xC000000000000000$
x86 Windows	$KPA = KVA \& 0x1FFFF000$
Mac OS X	$KPA = KVA$

As you can see the formula for Mac OS X, is the easiest existing formula.

### SMART TRANSLATION FORMULA

Using Quick Translation Formula, we can retrieve variables from `__DATA` section and initialized by `slave_pstart()` function of Mac OS X Kernel, which is called during the Operating System initialization.

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<sup>1</sup> 3.6 PAGING (VIRTUAL MEMORY) OVERVIEW.

There are ~4 variables which are interesting to perform the Smart Translation Formula: `IdlePDPT`, `IdlePDPT64`, `IdlePML4` and `IdlePTD`.

`IdlePML4` variable is initialized even on 32-bits Operating System. PML4 stands for Page Map Level 4 paging structure. This method can be used to address up to  $2^{27}$  pages, which spans a linear address space of  $2^{48}$  bytes.

Then, using `IdlePML4` variable we can cover a translation mechanism for a linear address space of  $2^{48}$  bytes even if the processor cannot do it. Internally, in kernel structures, Mac OS X is using 64-bits addressing for memory objects.

These variable are used later to initialize `kernel_map` and `kernel_pmap` kernel structures/variables.

Here is a common output of these variables under Mac OS X Leopard.

```
*_IdlePML4: [0x004EB00C] = 0x01219000
0x01219000: 27 A0 21 01 00 00 00 00 - 00 00 00 00 00 00 00 00 .....
0x01219010: 00 00 00 00 00 00 00 00 - 00 00 00 00 00 00 00 00 .....
0x01219020: 00 00 00 00 00 00 00 00 - 00 00 00 00 00 00 00 00 .....
0x01219030: 00 00 00 00 00 00 00 00 - 00 00 00 00 00 00 00 00 .....
0x01219040: 00 00 00 00 00 00 00 00 - 00 00 00 00 00 00 00 00 .....
0x01219050: 00 00 00 00 00 00 00 00 - 00 00 00 00 00 00 00 00 .....

*_IdlePDPT64: [0x004EB010] = 0x0121A000
0x0121A000: 27 C0 21 01 00 00 00 00 - 27 D0 21 01 00 00 00 00 .....
0x0121A010: 27 E0 21 01 00 00 00 00 - 27 F0 21 01 00 00 00 00 .....
0x0121A020: 00 00 00 00 00 00 00 00 - 00 00 00 00 00 00 00 00 .....
0x0121A030: 00 00 00 00 00 00 00 00 - 00 00 00 00 00 00 00 00 .....
0x0121A040: 00 00 00 00 00 00 00 00 - 00 00 00 00 00 00 00 00 .....
0x0121A050: 00 00 00 00 00 00 00 00 - 00 00 00 00 00 00 00 00 .....

*_IdlePDPT: [0x004EB008] = 0x0121B000
0x0121B000: 01 C0 21 01 00 00 00 00 - 01 D0 21 01 00 00 00 00 .....
0x0121B010: 01 E0 21 01 00 00 00 00 - 01 F0 21 01 00 00 00 00 .....
0x0121B020: 00 00 00 00 00 00 00 00 - 00 00 00 00 00 00 00 00 .....
0x0121B030: 00 00 00 00 00 00 00 00 - 00 00 00 00 00 00 00 00 .....
0x0121B040: 00 00 00 00 00 00 00 00 - 00 00 00 00 00 00 00 00 .....
0x0121B050: 00 00 00 00 00 00 00 00 - 00 00 00 00 00 00 00 00 .....

*_IdlePTD: [0x004EB004] = 0x0121C000
0x0121C000: 63 50 02 01 00 00 00 00 - 63 60 02 01 00 00 00 00 cP.....c`.....
0x0121C010: 63 70 02 01 00 00 00 00 - 63 80 02 01 00 00 00 00 cp.....c.....
0x0121C020: 23 90 02 01 00 00 00 00 - 23 A0 02 01 00 00 00 00 .....
0x0121C030: 63 B0 02 01 00 00 00 00 - 63 C0 02 01 00 00 00 00 c.....c.....
0x0121C040: 63 D0 02 01 00 00 00 00 - 63 E0 02 01 00 00 00 00 c.....c.....
0x0121C050: 63 F0 02 01 00 00 00 00 - 63 00 03 01 00 00 00 00 c.....c.....
```

## SYMBOLS

Symbols are a key element of volatile memory forensics without them an advanced analysis is impossible. Symbols of Microsoft Windows are available on a remote server as standalone files, but on Mac OS X symbols are directly stored inside the executable in a segment/section called `__LINKEDIT`.

The easiest way to retrieve kernel symbols is to extract them from the kernel executable of the hard-drive.

Symbols are firstly used to retrieve the address of memory variable for Smart Translation Formula.

## FAT HEADER

Mac OS X file format follows the FAT file format which contains magic signature of the header and the number of different architectures entries (i386, PowerPC or Both) inside the executable in big endian.

```
#define FAT_MAGIC 0xBEBABEFA

typedef struct _FAT_HEADER
{
    ULONG magic;
    ULONG nfat_arch;
} FAT_HEADER, *PFAT_HEADER;
```

To jump to the first architecture entry we add `sizeof(FAT_HEADER)` bytes to the pointer of the file header. Each entry uses the following definition, and also uses the big endian endianness.

```
typedef struct _FAT_ARCH
{
    cpu_type_t cputype;
    cpu_subtype_t cpusubtype;
    ULONG offset;
    ULONG size;
    ULONG align;
} FAT_ARCH, *PFAT_ARCH;
```

The first field, `cpu_type`, indicates to the loader what kind of architecture this entry defines using the following description:

```
typedef enum
{
    CPU_TYPE_VAX = 1,
    CPU_TYPE_ROMP = 2,
    CPU_TYPE_NS32032 = 4,
    CPU_TYPE_NS32332 = 5,
    CPU_TYPE_MC680x0 = 6,
    CPU_TYPE_I386 = 7,
    CPU_TYPE_MIPS = 8,
    CPU_TYPE_NS32532 = 9,
    CPU_TYPE_MC98000 = 10,
    CPU_TYPE_HPPA = 11,
    CPU_TYPE_ARM = 12,
    CPU_TYPE_MC88000 = 13,
    CPU_TYPE_SPARC = 14,
    CPU_TYPE_I860 = 15,
    CPU_TYPE_ALPHA = 16,
    CPU_TYPE_POWERPC = 18,
    /* APPLE LOCAL 64-bit */
    CPU_TYPE_POWERPC_64 = (18 | CPU_IS64BIT),
    /* APPLE LOCAL x86_64 */
    CPU_TYPE_X86_64 = (CPU_TYPE_I386 | CPU_IS64BIT)
} cpu_type_t;
```

And the third field, `offset`, contains the raw offset of the architecture header.

We assume index `x` is the id of the `CPU_TYPE_I386` architecture. So we have `FAT_ARCH[x].cputype` equals to `CPU_TYPE_I386` and `FAT_ARCH[x].offset` as new pointer offset to the `MACH_HEADER` structure.

## MACH HEADER

Now we have a pointer the i386 architecture binary using the following header definition and little-endian endianness.

```
#define MH_MAGIC 0xfeedface

typedef struct _MACH_HEADER
{
    ULONG Magic;
    cpu_type_t cputype;
    cpu_subtype_t cpusubtype;
    ULONG filetype;
    ULONG ncmds;
    ULONG sizeofcmds;
    ULONG flags;
} MACH_HEADER, *PMACH_HEADER;
```

This architecture validity can be verified using the `0xfeedface` magic key.

Now we can read what Apple calls *commands*, the field `MACH_HEADER.ncmds` indicates the number of commands inside the Mach-O binary.

We have to add `sizeof(MACH_HEADER)` to the Mach-O header pointer to have a pointer to the first command entry. There are different commands types and size of commands depends of their type. Most important commands types are `LC_SEGMENT` and `LC_SYMTAB`.

```
#define LC_SEGMENT 0x1 /* file segment to be mapped */
#define LC_SYMTAB 0x2 /* link-edit stab symbol table info (obsolete) */
```

And very two first fields contains information about the command's type and its size, using the following scheme:

```
typedef struct _LOAD_COMMAND {
    ULONG cmd; /* type of load command */
    ULONG cmdsize; /* total size of command in bytes */
} LOAD_COMMAND, *PLOAD_COMMAND;
```

Command type called `LC_SYMTAB`, contains raw pointers to two different tables. One, called `symoff`, with `NLIST` structures-based entries, and another, called `stroff`, with functions and variables names of each corresponding entry in the same order.

```
typedef struct _SYMTAB_COMMAND
{
    ULONG cmd;
    ULONG cmdsize;
```

```

        ULONG symoff;
        ULONG nsyms;
        ULONG stroff;
        ULONG strsize;
    } SYMTAB_COMMAND, *PSYMTAB_COMMAND;

typedef struct _NLIST
{
    ULONG n_strx;
    UCHAR n_type;
    UCHAR n_sect;
    USHORT n_desc;
    ULONG n_value;
} NLIST, *PNLIST;

```

Both `symoff` and `stroff` are pointer into the `__LINKEDIT` segment. Please note we have to add `FAT_ARCH[x].offset` value to these fields. And `n_value` field from `NLIST` structure contains the symbol offset.

Here is a short dump of symbols retrieved from Mac OS X Leopard kernel.

```

[000000] .constructors_used                0x0050F254
[000001] .destructors_used                   0x0050F25C
[000002] _AARPwakeup                         0x0029F6BD
[000003] _APTD                               0xFF7F8000
[000004] _APTDpde                            0xFFFF7FC0
[000005] _APTmap                              0xFF000000
[000006] _ASPgetmsg                          0x00293448
[000007] _ASPputmsg                          0x00293A3F
[000008] _ATPgetreq                          0x002A5A12
[000009] _ATPgetrsp                          0x002A5A79
[000010] _ATPsndreq                          0x002A592C
[000011] _ATPsndrsp                          0x002A599F
[000012] _ATgetmsg                           0x002A5844
[000013] _ATputmsg                           0x002A58B8
[... ]
[000223] _IOZeroTvalspec                     0x0050EE18
[000224] _IS_64BIT_PROCESS                   0x00373952
[000225] _IdlePDPT                          0x004EB008
[000226] _IdlePDPT64                        0x004EB010
[000227] _IdlePML4                          0x004EB00C
[000228] _IdlePTD                            0x004EB004
[000229] _InitGlobals                        0x002A3A08
[000230] _InsertKeyRecord                    0x003477C8
[000231] _InsertOffset                       0x003475A1
[000232] _InsertRecord                       0x00347703
[... ]
[002577] __ZN32IOServiceMessageUserNotificationC2EPK11OSMetaClass 0x004354AA
[002578] __ZN32IOServiceMessageUserNotificationC2Ev              0x0043560E
[002579] __ZN32IOServiceMessageUserNotificationD0Ev            0x0043553A
[002580] __ZN32IOServiceMessageUserNotificationD1Ev            0x0043551A
[002581] __ZN32IOServiceMessageUserNotificationD2Ev            0x004354FA
[002582] __ZN5IOCPU10gMetaClassE                               0x0052E838
[002583] __ZN5IOCPU10superClassE                               0x004BDE7C
[002584] __ZN5IOCPU11getCPUGroupEv                             0x004300A2
[002585] __ZN5IOCPU11getCPUStateEv                             0x00430082
[002586] __ZN5IOCPU11setCPUStateEm                             0x0043008E
[002587] __ZN5IOCPU11setPropertyEPK8OSSymbolP8OSObject         0x0042FE68
[... ]

```

[013859] _zt_ent_zindex	0x00299028
[013860] _zt_find_zname	0x00298D6D
[013861] _zt_getNextZone	0x0029910E
[013862] _zt_get_zmcast	0x00298F1D
[013863] _zt_remove_zones	0x00298CDD
[013864] _zt_set_zmap	0x002990B6

## INFORMATION EXTRACTION (ALSO KNOW AS ANALYSIS)

Once memory manager is functional, we can now proceed to the extraction of information such as process list and so on.

### MACHINE INFORMATION

Machine identification is a very important part to validate result. This section covers how to retrieve Darwin version, compilation date, number of CPUs and available memory on the current system.

There is a global variable, accessible from symbols, called `version` which contains a 100 bytes string with O.S. Type, O.S. Release version, username who compiled it.

There is another global variable, accessible from symbols, called `machine_info` defined by `machine_info` structure which contains information about CPUs and Memory of the target machine.

Definition of `machine_info` structure can be retrieved in `xnu/osfmk/mach/machine.h` header file.

Below is the definition of `machine_info` structure under Mac OS X Snow Leopard.

```
struct machine_info {
    integer_t major_version; /* kernel major version id */
    integer_t minor_version; /* kernel minor version id */
    integer_t max_cpus; /* max number of CPUs possible */
    uint32_t memory_size; /* size of memory in bytes, capped at 2 GB */
    uint64_t max_mem; /* actual size of physical memory */
    uint32_t physical_cpu; /* number of physical CPUs now available */
    integer_t physical_cpu_max; /* max number of physical CPUs possible */
    uint32_t logical_cpu; /* number of logical cpu now available */
    integer_t logical_cpu_max; /* max number of physical CPUs possible */
};
```

```
Darwin Kernel Version 9.0.0: Tue Oct 9 21:35:55 PDT 2007; root:xnu-1228~1/RELEASE_I386
Major version:      9
Minor version:     0
Max number of CPUs: 4
Size of physical memory: 1024 MB
Number of physical CPUs: 0
Number of logical CPUs: 1
```

Above is a screenshot of extraction information showing the target machine is running Mac OS X Leopard 10.5.0 with 1GB of physical memory.

### MOUNTED FILE SYSTEMS

Mounted file systems are defined by a global list-head, accessible from symbols, called `mountlist`. `mountlist` is a single link-list and contains a pointer called `next` which is a pointer to the next mounted file system entry both are defined by `mount` structure.

This structure contains 3 important fields including: file system type (`f_fstypename`), directory on which mounted (`f_mntonname`) and mounted file system (`f_mntfromname`).

Definition of `mount` structure can be retrieved in `xnu/bsd/sys/mount_internal.h` header file.

Below is the definition of `mount` structure under Mac OS X Snow Leopard.

```
/*
 * Structure per mounted file system. Each mounted file system has an
 * array of operations and an instance record. The file systems are
 * put on a doubly linked list.
 */

struct mount {
    TAILQ_ENTRY(mount) mnt_list;          /* mount list */
    int32_t mnt_count;                   /* reference on the mount */
    lck_mtx_t mnt_mlock;                 /* mutex that protects mount point */
    struct vfsops *mnt_op;                /* operations on fs */
    struct vfstable *mnt_vtable;         /* configuration info */
    struct vnode *mnt_vnodecovered;      /* vnode we mounted on */
    struct vnodelst mnt_vnodelist;       /* list of vnodes this mount */
    struct vnodelst mnt_workerqueue;     /* list of vnodes this mount */
    struct vnodelst mnt_newvnodes;      /* list of vnodes this mount */
    uint32_t mnt_flag;                   /* flags */
    uint32_t mnt_kern_flag;              /* kernel only flags */
    uint32_t mnt_lflag;                  /* mount life cycle flags */
    uint32_t mnt_maxsymlinklen;         /* max size of short symlink */
    struct vfsstatfs mnt_vfsstat;        /* cache of filesystem stats */
    qaddr_t mnt_data;                   /* private data */
    /* Cached values of the IO constraints for the device */
    uint32_t mnt_maxreadcnt;             /* Max. byte count for read */
    uint32_t mnt_maxwritecnt;           /* Max. byte count for write */
    uint32_t mnt_segreadcnt;            /* Max. segment count for read */
    uint32_t mnt_segwritecnt;           /* Max. segment count for write */
    uint32_t mnt_maxsegreadsize;        /* Max. segment read size */
    uint32_t mnt_maxsegwritesize;       /* Max. segment write size */
    uint32_t mnt_alignmentmask;         /* Mask of bits that aren't addressable
via DMA */
    uint32_t mnt_devblocksize;          /* the underlying device block size */
    uint32_t mnt_ioqueue_depth;         /* the maximum number of commands a
device can accept */
    uint32_t mnt_ioscale;               /* scale the various throttles/limits imposed
on the amount of I/O in flight */
    uint32_t mnt_ioflags;               /* flags for underlying device */
    pending_io_t mnt_pending_write_size; /* byte count of pending writes */
    pending_io_t mnt_pending_read_size; /* byte count of pending reads */

    lck_rw_t mnt_rwlock;                /* mutex readwrite lock */
    lck_mtx_t mnt_renamelock;           /* mutex that serializes renames that change
shape of tree */
    vnode_t mnt_devvp;                 /* the device mounted on for local file systems */
};
```

```

uint32_t    mnt_devbsdunit; /* the BSD unit number of the device */
void *mnt_throttle_info; /* used by the throttle code */
int32_t mnt_crossref; /* refernces to cover lookups crossing into mp
*/
int32_tmnt_iterref; /* refernces to cover iterations; drained makes it
-ve */

/* XXX 3762912 hack to support HFS filesystem 'owner' */
uid_t      mnt_fsowner;
gid_t      mnt_fsgroup;

struct label    *mnt_mntlabel;          /* MAC mount label */
struct label    *mnt_fslabel;          /* MAC default fs label */

/*
 * cache the rootvp of the last mount point
 * in the chain in the mount struct pointed
 * to by the vnode sitting in '/'
 * this cache is used to shortcircuit the
 * mount chain traversal and allows us
 * to traverse to the true underlying rootvp
 * in 1 easy step inside of 'cache_lookup_path'
 *
 * make sure to validate against the cached vid
 * in case the rootvp gets stolen away since
 * we don't take an explicit long term reference
 * on it when we mount it
 */
vnode_t      mnt_realrootvp;
uint32_t     mnt_realrootvp_vid;
/*
 * bumped each time a mount or unmount
 * occurs... its used to invalidate
 * 'mnt_realrootvp' from the cache
 */
uint32_t     mnt_generation;
/*
 * if 'MNTK_AUTH_CACHE_TIMEOUT' is
 * set, then 'mnt_authcache_ttl' is
 * the time-to-live for the per-vnode authentication cache
 * on this mount... if zero, no cache is maintained...
 * if 'MNTK_AUTH_CACHE_TIMEOUT' isn't set, its the
 * time-to-live for the cached lookup right for
 * volumes marked 'MNTK_AUTH_OPAQUE'.
 */
int          mnt_authcache_ttl;
/*
 * The proc structure pointer and process ID form a
 * sufficiently unique duple identifying the process
 * hosting this mount point. Set by vfs_markdependency()
 * and utilized in new_vnode() to avoid reclaiming vnodes
 * with this dependency (radar 5192010).
 */
pid_t       mnt_dependent_pid;
void        *mnt_dependent_process;
};

```

id#	type	mounted on	mounted from
0	hfs	/	nfo
1	devfs	/dev	devfs
2	fdesc	/dev	fdesc
3	autofs	/net	map -hosts
4	autofs	/home	map auto_home
5	hfs	/Volumes/VMware Tools	né
6	hfs	/Volumes/OSXBAK	/dev/disk2s1
7	msdos	/Volumes/FATBACK	/dev/disk2s2

Above is a screenshot of mounted file systems including an external hard-drive.

## BSD PROCESSES

Every Operating System uses user-land processes, it is one of the key element of a working O.S.

Loaded processes are stored into `proc` structure which contains a double-list to walk into the list. There is a global variable, retrievable from symbols, called `kernproc` is the list-head of BSD processes list.

`p_list` field is a double link-list which contains a pointer to both, the previous and the next process.

Definition of `proc` structure can be retrieved in `xnu/bsd/sys/proc_internal.h` header file.

Below is the definition of `proc` structure under Mac OS X Snow Leopard.

```

/*
 * Description of a process.
 *
 * This structure contains the information needed to manage a thread of
 * control, known in UN*X as a process; it has references to substructures
 * containing descriptions of things that the process uses, but may share
 * with related processes. The process structure and the substructures
 * are always addressible except for those marked "(PROC ONLY)" below,
 * which might be addressible only on a processor on which the process
 * is running.
 */
struct proc {
    LIST_ENTRY(proc) p_list;          /* List of all processes. */

    pid_t          p_pid;             /* Process identifier. (static)*/
    void *         task;              /* corresponding task (static)*/
    struct proc *  p_pptr;            /* Pointer to parent process. (LL) */
    pid_t          p_ppid;            /* process's parent pid number */
    pid_t          p_pgrpid;          /* process group id of the process (LL)*/

    lck_mtx_t      p_mlock;          /* mutex lock for proc */

    char           p_stat;             /* S* process status. (PL)*/
    char           p_shutdownstate;
    char           p_kdebug;          /* P_KDEBUG eq (CC)*/
    char           p_btrace;          /* P_BTRACE eq (CC)*/

    LIST_ENTRY(proc) p_pgl;           /* List of processes in pgrp. (PGL) */
    LIST_ENTRY(proc) p_sibling;       /* List of sibling processes. (LL)*/
    LIST_HEAD(, proc) p_children;     /* Pointer to list of children. (LL)*/
    TAILQ_HEAD(, uthread) p_uthlist; /* List of utthreads (PL) */

```

```

LIST_ENTRY(proc) p_hash;          /* Hash chain. (LL)*/
TAILQ_HEAD( ,eventqelt) p_evlist; /* (PL) */

lck_mtx_t    p_fdmlock; /* proc lock to protect fdesc */

/* substructures: */
kauth_cred_t    p_ucred; /* Process owner's identity. (PL) */
struct filedesc *p_fd; /* Ptr to open files structure. (PFDL) */
struct pstats *p_stats; /* Accounting/statistics (PL). */
struct plimit *p_limit; /* Process limits.(PL) */

struct sigacts *p_sigacts; /* Signal actions, state (PL) */
int p_siglist; /* signals captured back from threads */
lck_spin_t p_slock; /* spin lock for itimer/profil protection */

#define p_rlimit p_limit->pl_rlimit

struct plimit *p_olimit; /* old process limits - not inherited by
child (PL) */
unsigned int p_flag; /* P_* flags. (atomic bit ops) */
unsigned int p_lflag; /* local flags (PL) */
unsigned int p_listflag; /* list flags (LL) */
unsigned int p_ladvflag; /* local adv flags (atomic) */
int p_refcount; /* number of outstanding users(LL) */
int p_childrencnt; /* children holding ref on parent (LL) */
int p_parentref; /* children lookup ref on parent (LL) */

pid_t p_oppid; /* Save parent pid during ptrace. XXX */
u_int p_xstat; /* Exit status for wait; also stop signal. */

#ifdef _PROC_HAS_SCHEDINFO_
/* may need cleanup, not used */
u_int p_estcpu; /* Time averaged value of p_cpticks.(used by aio and
proc_comapre) */
fixpt_t p_pctcpu; /* %cpu for this process during p_swtime (used by
aio)*/
u_int p_slptime; /* used by proc_compare */
#endif /* _PROC_HAS_SCHEDINFO_ */

struct itimerval p_realtimer; /* Alarm timer. (PSL) */
struct timeval p_rtime; /* Real time.(PSL) */
struct itimerval p_vtimer_user; /* Virtual timers.(PSL) */
struct itimerval p_vtimer_prof; /* (PSL) */

struct timeval p_rlim_cpu; /* Remaining rlim cpu value.(PSL) */
int p_debugger; /* NU 1: can exec set-bit programs if suser */
boolean_t sigwait; /* indication to suspend (PL) */
void *sigwait_thread; /* 'thread' holding sigwait(PL) */
void *exit_thread; /* Which thread is exiting(PL) */
int p_vforkcnt; /* number of outstanding vforks(PL) */
void * p_vforkact; /* activation running this vfork proc)(static) */
int p_fpdrainwait; /* (PFDL) */
pid_t p_contproc; /* last PID to send us a SIGCONT (PL) */

/* Following fields are info from SIGCHLD (PL) */
pid_t si_pid; /* (PL) */

```

```

u_int    si_status; /* (PL) */
u_int    si_code;   /* (PL) */
uid_t    si_uid;    /* (PL) */

void * vm_shm; /* (SYSV SHM Lock) for sysV shared memory */

#if CONFIG_DTRACE
user_addr_t p_dtrace_argv; /* (write once, read only after that) */
user_addr_t p_dtrace_envp; /* (write once, read only after that) */
lck_mtx_t   p_dtrace_sprlock; /* sun proc lock emulation */
int         p_dtrace_probes; /* (PL) are there probes for this proc? */
u_int       p_dtrace_count; /* (sprlock) number of DTrace tracepoints */
struct      dtrace_ptss_page* p_dtrace_ptss_pages; /* (sprlock) list of user
ptss pages */
struct      dtrace_ptss_page_entry* p_dtrace_ptss_free_list; /* (atomic)
list of individual ptss entries */
struct      dtrace_helpers* p_dtrace_helpers; /* (dtrace_lock) DTrace per-
proc private */
struct      dof_ioctl_data* p_dtrace_lazy_dofs; /* (sprlock) unloaded
dof_helper_t's */
#endif /* CONFIG_DTRACE */

/* XXXXXXXXXXXXXXXX BCOPY'ed on fork XXXXXXXXXXXXXXXXXXXX */
/* The following fields are all copied upon creation in fork. */
#define      p_startcopy p_arghlen

u_int       p_arghlen; /* Length of process arguments. */
int         p_argc; /* saved argc for sysctl_procargs() */
user_addr_t user_stack; /* where user stack was allocated */
struct      vnode *p_textvp; /* Vnode of executable. */
off_t       p_textoff; /* offset in executable vnode */

sigset_t    p_sigmask; /* DEPRECATED */
sigset_t    p_sigignore; /* Signals being ignored. (PL) */
sigset_t    p_sigcatch; /* Signals being caught by user. (PL) */

u_char      p_priority; /* (NU) Process priority. */
u_char      p_resv0; /* (NU) User-priority based on p_cpu and
p_nice. */
char        p_nice; /* Process "nice" value. (PL) */
u_char      p_resv1; /* (NU) User-priority based on p_cpu and p_nice. */

#if CONFIG_MACF
int         p_mac_enforce; /* MAC policy enforcement control */
#endif

char        p_comm[MAXCOMLEN+1];
char        p_name[(2*MAXCOMLEN)+1]; /* PL */

struct      pgrp *p_pgrp; /* Pointer to process group. (LL) */
int         p_iopol_disk; /* disk I/O policy (PL) */
uint32_t    p_csflags; /* flags for codesign (PL) */
uint32_t    p_paction; /* action for process control on starvation */
uint8_t     p_uuid[16]; /* from LC_UUID load command */

/* End area that is copied on creation. */
/* XXXXXXXXXXXXXXXX End of BCOPY'ed on fork (AIOLock)XXXXXXXXXXXXXXXXX */

```

```

#define      p_endcopy      p_aio_total_count
      int      p_aio_total_count; /* all allocated AIO requests for this
proc */
      int      p_aio_active_count; /* all unfinished AIO requests for this
proc */
      TAILQ_HEAD( , aio_workq_entry ) p_aio_activeq; /* active async IO
requests */
      TAILQ_HEAD( , aio_workq_entry ) p_aio_doneq; /* completed async IO
requests */

      struct klist p_klist; /* knote list (PL ?)*/

      struct      rusage *p_ru; /* Exit information. (PL) */
      thread_t      p_signalholder;
      thread_t      p_transholder;

      /* DEPRECATE following field */
      u_short      p_acflag; /* Accounting flags. */

      struct lctx *p_lctx; /* Pointer to login context. */
      LIST_ENTRY(proc) p_lclist; /* List of processes in lctx. */
      user_addr_t      p_threadstart; /* pthread start fn */
      user_addr_t      p_wqthread; /* pthread workqueue fn */
      int      p_pthsize; /* pthread size */
      user_addr_t p_targconc; /* target concurrency ptr */
      void *      p_wqptr; /* workq ptr */
      int      p_wqsize; /* allocated size */
      boolean_t      p_wqiniting; /* semaphore to serialize wq_open */
      lck_spin_t      p_wqlck; /* lock to protect work queue */
      struct timeval p_start; /* starting time */
      void *      p_rcall;
      int      p_ractive;
      int      p_idversion; /* version of process identity */
      void *      p_pthhash; /* pthread waitqueue hash */
#if DIAGNOSTIC
      unsigned int p_fdlock_pc[4];
      unsigned int p_fdunlock_pc[4];
#endif
#if SIGNAL_DEBUG
      unsigned int lockpc[8];
      unsigned int unlockpc[8];
#endif /* SIGNAL_DEBUG */
#endif /* DIAGNOSTIC */
      uint64_t      p_dispatchqueue_offset;
};

```

Pointer to the process group, `pgrp` structure, allows us to retrieve the username of the person who launched the program because this structure contains a pointer to a structure called `session` with the username.

Definition of `pgrp` structure can also be retrieved in `xnu/bsd/sys/proc_internal.h` header file.

Below is the definition of `pgrp` structure under Mac OS X Snow Leopard.

```

/*
 * One structure allocated per process group.
 */
struct pgrp {

```

```

LIST_ENTRY(pgrp) pg_hash; /* Hash chain. (LL) */
LIST_HEAD(, proc) pg_members; /* Pointer to pgrp members. (PGL) */
struct session *pg_session; /* Pointer to session. (LL) */
pid_t pg_id; /* Pgrp id. (static) */
int pg_jobc; /* # procs qualifying pgrp for job control (PGL) */
int pg_membercnt; /* Number of processes in the pprocess group (PGL) */
int pg_refcount; /* number of current iterators (LL) */
unsigned int pg_listflags; /* (LL) */
lck_mtx_t pg_mlock; /* mutex lock to protect pgrp */
};

```

Definition of session structure can also be retrieved in xnu/bsd/sys/proc\_internal.h header file.

Below is the definition of session structure under Mac OS X Snow Leopard.

```

/*
 * One structure allocated per session.
 */
struct session {
    int s_count; /* Ref cnt; pgrps in session. (LL) */
    struct proc *s_leader; /* Session leader.(static) */
    struct vnode *s_ttyp; /* Vnode of controlling terminal.(SL) */
    int s_ttyvid; /* Vnode id of the controlling terminal (SL) */
    struct tty *s_ttyp; /* Controlling terminal. (SL + ttyp != NULL) */
    pid_t s_ttypgrpid; /* tty's pgrp id */
    pid_t s_sid; /* Session ID (static) */
    char s_login[MAXLOGNAME]; /* Setlogin() name.(SL) */
    int s_flags; /* Session flags (s_mlock) */
    LIST_ENTRY(session) s_hash; /* Hash chain.(LL) */
    lck_mtx_t s_mlock; /* mutex lock to protect session */
    int s_listflags;
};

```

s\_login field contains the name of the username in ASCII.

task#	pid	parent pid	name	username	started time
1	0	0	kernel_task		Thu 2009-March-26 12:44:43 (W. Europe Standard Time)
2	1	0	launchd	nfinfi	Thu 2009-March-26 12:44:43 (W. Europe Standard Time)
3	10	1	kextd	root	Thu 2009-March-26 12:44:45 (W. Europe Standard Time)
4	11	1	notified	root	Thu 2009-March-26 12:44:45 (W. Europe Standard Time)
5	12	1	syslogd	root	Thu 2009-March-26 12:44:46 (W. Europe Standard Time)
6	14	1	ntpd	root	Thu 2009-March-26 12:44:47 (W. Europe Standard Time)
7	16	1	update	root	Thu 2009-March-26 12:44:47 (W. Europe Standard Time)
8	19	1	securityd	root	Thu 2009-March-26 12:44:47 (W. Europe Standard Time)
9	21	1	nds	root	Thu 2009-March-26 12:44:47 (W. Europe Standard Time)
10	22	1	ndNSResponder		Thu 1970-January-01 01:00:00 (W. Europe Standard Time)
11	23	1	loginwindow	nfinfi	Thu 2009-March-26 12:44:47 (W. Europe Standard Time)
12	24	1	KernelEventAgent	root	Thu 2009-March-26 12:44:47 (W. Europe Standard Time)
13	26	1	hidd	root	Thu 2009-March-26 12:44:47 (W. Europe Standard Time)
14	27	1	fsevents	root	Thu 2009-March-26 12:44:47 (W. Europe Standard Time)
15	28	1	dynamic_pager	root	Thu 2009-March-26 12:44:47 (W. Europe Standard Time)
16	31	1	diskarbitrationd	root	Thu 2009-March-26 12:44:47 (W. Europe Standard Time)
17	32	1	DirectoryService	root	Thu 2009-March-26 12:44:47 (W. Europe Standard Time)
18	34	1	configd	root	Thu 2009-March-26 12:44:47 (W. Europe Standard Time)
19	37	1	autofs	root	Thu 2009-March-26 12:44:47 (W. Europe Standard Time)
20	38	1	socketfilterfw	root	Thu 2009-March-26 12:44:47 (W. Europe Standard Time)
21	40	1			Thu 1970-January-01 01:00:00 (W. Europe Standard Time)
22	46	1	coreservicesd	_securityagent	Thu 2009-March-26 12:44:51 (W. Europe Standard Time)
23	48	1	WindowServer	root	Thu 2009-March-26 12:44:51 (W. Europe Standard Time)
24	59	1	launchd	nfinfi	Thu 2009-March-26 12:44:53 (W. Europe Standard Time)
25	71	1	coreaudiod	nobody	Thu 2009-March-26 12:45:03 (W. Europe Standard Time)
26	78	59	Spotlight	nfinfi	Thu 2009-March-26 12:45:04 (W. Europe Standard Time)
27	79	59	UserEventAgent	nfinfi	Thu 2009-March-26 12:45:04 (W. Europe Standard Time)
28	80	59	Dock	nfinfi	Thu 2009-March-26 12:45:04 (W. Europe Standard Time)
29	81	59	SystemUIServer	nfinfi	Thu 2009-March-26 12:45:04 (W. Europe Standard Time)
30	82	59	Finder	nfinfi	Thu 2009-March-26 12:45:04 (W. Europe Standard Time)
31	83	59	AISServer	nfinfi	Thu 2009-March-26 12:45:04 (W. Europe Standard Time)
32	85	59	pboard	nfinfi	Thu 2009-March-26 12:45:04 (W. Europe Standard Time)
33	96	14	ntpd	root	Thu 2009-March-26 12:46:10 (W. Europe Standard Time)
34	97	59	Terminal	nfinfi	Thu 2009-March-26 12:50:47 (W. Europe Standard Time)
35	98	97	login	nfinfi	Thu 2009-March-26 12:50:47 (W. Europe Standard Time)
36	99	98	bash	nfinfi	Thu 2009-March-26 12:50:48 (W. Europe Standard Time)
37	128	59	Preview	nfinfi	Thu 2009-March-26 12:56:36 (W. Europe Standard Time)
38	211	59	Xcode	nfinfi	Thu 2009-March-26 12:59:32 (W. Europe Standard Time)
39	228	97	login	nfinfi	Thu 2009-March-26 13:30:08 (W. Europe Standard Time)

Above is a sample screenshot of a processes list, mainly executed by nfinfi user around March 2009.

## KERNEL EXTENSIONS (ALSO KNOWN AS DRIVERS, KERNEL MODULES)

Kernel-Mode, the God Mode, is the most privileged level of an Operating System. Loaded Kernel Extensions can be retrieved by a global list-head variable, accessible from symbols, called `kmod` defined by `kmod_info` structure.

`next` field points to the next kernel extension.

Definition of `session` structure can also be retrieved in `xnu/osfmk/mach/kmod.h` header file.

Below is the definition of `kmod_info` structure under Mac OS X Snow Leopard.

```
typedef struct kmod_info {
    struct kmod_info *next;
    int info_version; // version of this structure
    int id;
    char name[KMOD_MAX_NAME];
    char version[KMOD_MAX_NAME];
    int reference_count; // # refs to this
    kmod_reference_t *reference_list; // who this refs
    vm_address_t address; // starting address
    vm_size_t size; // total size
    vm_size_t hdr_size; // unwired hdr size
    kmod_start_func_t *start;
    kmod_stop_func_t *stop;
} kmod_info_t;
```

As you can see here we have both kernel extensions image base start and size. Since we have a functional kernel address space, we can easily extract the image of the kernel extension.

id	ref	address	size	code size	name (version)
106	0	0x226F0000	0x00003000	0x00002000	con.apple.driver.iTunesPhoneDriver (1.0)
104	0	0x226D0000	0x0000A000	0x00009000	con.apple.iokit.IOUSBMassStorageClass (2.0.0)
102	0	0x226D1000	0x0001B000	0x0001A000	con.apple.filesystems.ntfs (2.0)
97	0	0x225ED000	0x0000C000	0x0000B000	con.apple.filesystems.msdosfs (1.5)
94	0	0x225CA000	0x00002000	0x00001000	con.osxbook.kext.KernelMemoryAccess (1.0.0)
93	0	0x00AC1000	0x00005000	0x00004000	con.apple.driver.AppleHWSensor (1.7.0d0)
92	0	0x00D58000	0x0000B000	0x0000A000	con.apple.filesystems.autofs (2.0.0d1)
91	0	0x00D35000	0x00023000	0x00022000	con.apple.driver.AppleHDA (1.4.0a23)
90	1	0x00CF4000	0x00041000	0x00040000	con.apple.driver.DspFuncLib (1.0.0a1)
89	0	0x00656000	0x00003000	0x00002000	con.apple.Dont_Steal_Mac_OS_X (6.0.2)
88	0	0x00DC6000	0x00007000	0x00006000	con.apple.iokit.CHUDUtils (200)
87	0	0x00C21000	0x0000C000	0x0000B000	con.apple.iokit.CHUDProf (207)
86	0	0x00DCD000	0x00082000	0x00081000	con.apple.GeForce (5.1.6)
85	2	0x00C1C000	0x00005000	0x00004000	con.apple.iokit.CHUDKernLib (196)
84	0	0x00909000	0x001B8000	0x001B7000	con.apple.nvidia.nv40hal (5.1.6)
83	2	0x006DC000	0x0022D000	0x0022C000	con.apple.NVDAResman (5.1.6)
82	2	0x006CE000	0x0000E000	0x0000D000	con.apple.iokit.IONDRVSupport (1.5)
81	3	0x006B2000	0x0001C000	0x0001B000	con.apple.iokit.IOGraphicsFamily (1.5)
80	0	0x00C90000	0x00007000	0x00006000	con.apple.driver.AppleHDAController (1.4.0a23)
79	2	0x00C8A000	0x00006000	0x00005000	con.apple.iokit.IOHDAFamily (1.4.0a23)
78	0	0x00C03000	0x00009000	0x00008000	con.apple.iokit.IOFireWireIP (1.7.0)
77	0	0x0069F000	0x00003000	0x00002000	con.apple.driver.AppleUSBDisplays (2.0)
76	0	0x00E4F000	0x00004000	0x00003000	con.apple.driver.AudioIPCDriver (1.0.4)
75	3	0x00CDD000	0x00017000	0x00016000	con.apple.iokit.IOAudioFamily (1.6.4b7)
74	1	0x00CDA000	0x00003000	0x00002000	con.apple.kext.OSvKernDSPLib (1.1)
73	0	0x00AC6000	0x0000A000	0x00009000	con.apple.driver.AppleMCEDriver (1.1.5b4)
72	0	0x00C71000	0x0000A000	0x00009000	con.apple.driver.ACPI_SMC_PlatformPlugin (3.0.0d11)
71	1	0x00C63000	0x0000E000	0x0000D000	con.apple.driver.IOPPlatformPluginFamily (3.0.0d11)
70	2	0x0064D000	0x00009000	0x00008000	con.apple.driver.AppleSMC (2.0.0d5)
69	0	0x00B0A000	0x00003000	0x00002000	con.apple.driver.AppleLPC (1.2.2)
68	0	0x00699000	0x00003000	0x00002000	con.apple.driver.AppleUSBHIDMouse (1.2.0b3)
67	1	0x00691000	0x00003000	0x00002000	con.apple.driver.AppleHIDMouse (1.2.0b3)
66	2	0x00694000	0x00005000	0x00004000	con.apple.iokit.IOUSBHIDDriver (3.0.3)
65	0	0x00647000	0x00002000	0x00001000	con.apple.driver.AppleUSBMergeNub (3.0.3)
64	0	0x00C4D000	0x00004000	0x00003000	con.apple.driver.AppleUSBComposite (3.0.0)
62	0	0x00BD4000	0x00016000	0x00015000	con.apple.iokit.IOSCSIMultimediaCommandsDevice (2.0.0)
61	2	0x00BBE000	0x00016000	0x00015000	con.apple.iokit.IOSCSIBlockCommandsDevice (2.0.0)
60	1	0x00BB9000	0x00005000	0x00004000	con.apple.iokit.IOBSDStorageFamily (1.5)
59	2	0x00BB3000	0x00006000	0x00005000	con.apple.iokit.IODVDStorageFamily (1.5)
58	3	0x00BA8000	0x00008000	0x00007000	con.apple.iokit.IOCDSStorageFamily (1.5)
57	0	0x00B04000	0x00006000	0x00005000	con.apple.iokit.SCSITaskUserClient (2.0.0)
56	0	0x00C97000	0x00005000	0x00004000	con.apple.driver.XsanFilter (2.7.91)

Above is a screenshot of a loaded kext lists.

## SYSTEM CALLS

The very first step is to localize the syscall table, called `sysent`, which is a non accessible variable from symbols. So using a magic trick we can retrieve its offset through `nsysent` exported variable which contains the number of syscall entries.

Under Mac OS X Leopard (10.5), as explained by Jesse D’Aguanno at BH US 2008, we have to add 0x20 to `nsysent` offset to obtain the offset of `sysent` table.

Under Mac OS X Snow Leopard (10.6), we have to proceed with a different methodology. First, we have to retrieve the value of `nsysent` variable, then we multiply its value with the size of `sysent` structure, and then we subtract this value to `nsysent` offset to obtain the offset of `sysent` table.

Definition of `sysent` structure can also be retrieved in `xnu/bsd/sys/sysent.h` header file.

Below is the definition of `sysent` structure under Mac OS X Snow Leopard.

```
struct sysent {
    /* system call table */
    int16_t sy_narg; /* number of args */
    int8_t sy_resv; /* reserved */
    int8_t sy_flags; /* flags */
    sy_call_t *sy_call; /* implementing function */
    sy_munge_t *sy_arg_munge32; /* system call arguments munger for 32-bit process */
};
```

```

sy_munge_t  *sy_arg_munge64; /* system call arguments munger for 64-bit
process */
int32_t sy_return_type; /* system call return types */
uint16_t sy_arg_bytes; /* Total size of arguments in bytes for
* 32-bit system calls
*/
};

```

id#	offset	name	table
0	0x003907F5	_nosys	[OK]
1	0x00376F34	_exit	[OK]
2	0x00378B4A	_fork	[OK]
3	0x00390CAE	_read	[OK]
4	0x0039134C	_write	[OK]
5	0x001E425C	_open	[OK]
6	0x0036C75E	_close	[OK]
7	0x00375EB2	_wait4	[OK]
8	0x003907F5	_nosys	[OK]
9	0x001E4932	_link	[OK]
10	0x001E5540	_unlink	[OK]
11	0x003907F5	_nosys	[OK]
12	0x001E3925	_chdir	[OK]
13	0x001E3723	_fchdir	[OK]
14	0x001E43E8	_mknod	[OK]
15	0x001E6FD1	_chmod	[OK]
16	0x001E74B7	_chown	[OK]
17	0x0037A52D	_obreak	[OK]
18	0x001E335E	_getfsstat	[OK]
19	0x003907F5	_nosys	[OK]
20	0x0037DE30	_getpid	[OK]
21	0x003907F5	_nosys	[OK]
22	0x003907F5	_nosys	[OK]
23	0x0037E92E	_setuid	[OK]
24	0x0037DF0D	_getuid	[OK]
25	0x0037DF21	_geteuid	[OK]
26	0x0038C823	_ptrace	[OK]
27	0x003B0A4E	_recvnsg	[OK]
28	0x003B1701	_sendnsg	[OK]
29	0x003B07D8	_recvfrom	[OK]
30	0x003AFE73	_accept	[OK]
31	0x003B0EC4	_getpeername	[OK]
32	0x003B0CDA	_getsockname	[OK]
33	0x001E5D2D	_access	[OK]
34	0x001E6BD7	_chflags	[OK]
35	0x001E6C88	_fchflags	[OK]
36	0x001E22B5	_sync	[OK]
37	0x003836B2	_kill	[OK]
38	0x003907F5	_nosys	[OK]
39	0x0037DE42	_getppid	[OK]
40	0x003907F5	_nosys	[OK]
41	0x0036E487	_dup	[OK]
42	0x00394912	_pipe	[OK]
43	0x0037DFC7	_getegid	[OK]
44	0x0038FBA6	_profil	[OK]
45	0x003907F5	_nosys	[OK]
46	0x00382075	_sigaction	[OK]
47	0x0037DFB3	_getgid	[OK]
48	0x003829F2	_sigprocmask	[OK]
49	0x0037E544	_getlogin	[OK]
50	0x0037E5E5	_setlogin	[OK]
51	0x003582A7	_acct	[OK]
52	0x00381125	_sigpending	[OK]
53	0x00381539	_sigaltstack	[OK]
54	0x0039160C	_ioctl	[OK]
55	0x0038C732	_reboot	[OK]
56	0x001E9F24	_revoke	[OK]
57	0x001E4E09	_symlink	[OK]
58	0x001E6923	_readlink	[OK]

Above is a picture showing a list of syscalls from `sysent` table.

Integrity checks are done if entry value does not give the function name value. It does not sound complicated but this trick was enough to detect Jesse D'Aguanno Rootkit presented at HAR2009.

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